

# From $1/\sqrt{n}$ to $1/n$ : Accelerating SDE Simulation with Cubature Formulae

Peter Koepernik<sup>1,\*</sup>   Thomas Coxon<sup>2,\*</sup>   James Foster<sup>3</sup>

<sup>1</sup>University of Oxford   <sup>2</sup>Loughborough University   <sup>3</sup>University of Bath

\*Equal contribution

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# Setting

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Consider an SDE

$$\begin{aligned} dy_t &= \mu(y_t) dt + \sigma(y_t) dW_t, \\ y_0 &= \xi, \end{aligned} \tag{1}$$

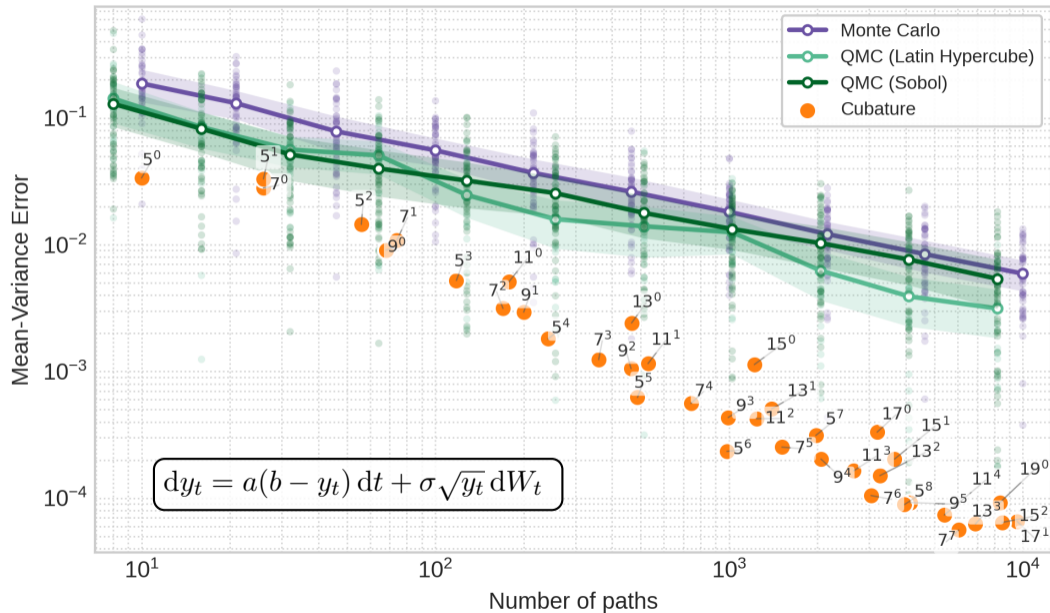
where  $y_t$  takes values in  $\mathbb{R}^e$ , and  $\mu, \sigma: \mathbb{R}^e \rightarrow \mathbb{R}^e$  are the *drift* and *diffusion*, respectively.

Statistics like  $\mathbb{E}[f(y_T)]$  for some  $f: \mathbb{R}^e \rightarrow \mathbb{R}$  are traditionally estimated with Monte-Carlo:

$$\mathbb{E}[f(y_T)] \approx \frac{1}{M} \sum_{i=1}^M f(y_T(\omega_i)) \quad \leftarrow \text{error } O(1/\sqrt{M})$$

Here  $y_T(\omega) \in \mathbb{R}^e$  denotes the solution at time  $T$  of (1) driven by a path  $\omega: [0, T] \rightarrow \mathbb{R}$ , and  $\omega_i$  are independent samples of Brownian motion.

# Cox-Ingersoll-Ross (CIR) Model



# Cubatures in Numerical Integration

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Numerical integration of a function  $f: [0, 1] \rightarrow \mathbb{R}$  can also be done with Monte-Carlo:

$$\int_0^1 f(x) dx = \mathbb{E}[f(U)] \approx \frac{1}{M} \sum_{i=1}^M f(u_i)$$

for  $U \sim \mathcal{U}([0, 1])$  and independent samples  $u_i \sim \mathcal{U}([0, 1])$ . The error is  $O(1/\sqrt{M})$ .

One can do much better with uniformly spaced points, which have error  $O(1/M)$ , and even better with *cubature formulae*: deterministic collections  $(\lambda_i, u_i)_{i=1}^M$  of points  $u_i \in [0, 1]$  with weights  $\lambda_i \geq 0$ ,  $\sum_i \lambda_i = 1$ , that are chosen in such a way that the quadrature approximation

$$\int_0^1 f(x) dx \approx \sum_{i=1}^M \lambda_i f(u_i)$$

is exact for polynomials  $f$  with degree up to some  $D \in \mathbb{N}$ , the *degree* of the cubature.

# Cubatures on Wiener Space

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Can we do something similar in the SDE setting?

$$\mathbb{E}[f(y_T)] \approx \sum_{i=1}^M \lambda_i f(y_T(\omega_i)),$$

where the i.i.d. paths  $\omega_i: [0, T] \rightarrow \mathbb{R}$  and uniform weights  $\lambda_i = 1/M$  are replaced by a *deterministic* collection  $(\lambda_i, \omega_i)_{i=1}^M$  of paths  $\omega_i$  and weights  $\lambda_i \geq 0$  that matches the expected Brownian “signature” up to some degree  $D$ :

$$\sum_{i=1}^M \lambda_i S_{0,T}^{(k)}(\omega_i) = \mathbb{E}[S_{0,T}^{(k)}(W)]$$

for all  $k \in \{0, \dots, D\}$ .

This is called a *Cubature on Wiener space* of degree  $D$  [Lyons and Victoir, 2004].

# Cubatures on Wiener Space

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The signature is made up of terms of the form

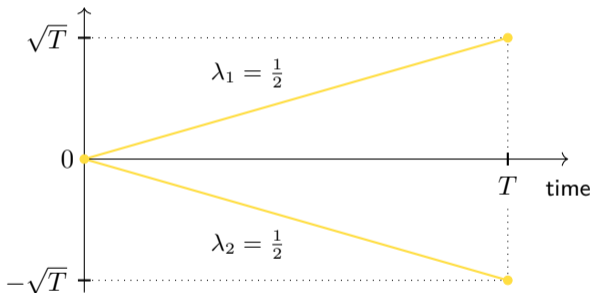
$$S_{0,T}^{(i_1 \dots i_k)}(\omega) = \int_{0 < t_1 < \dots < t_k < T} d\omega^{(i_1)}(t_1) \dots d\omega^{(i_k)}(t_k),$$

where  $\omega^{(0)}(t) = t$  and  $\omega^{(1)}(t) = \omega(t)$ . The *degree* of the term is the number of indices, where zeros are counted twice.

- $S^{(1)}(\omega) = \{\int_0^T d\omega_t\}$
- $S^{(2)}(\omega) = \{\int_0^T d\omega_{t_1} d\omega_{t_2}, \int_0^T dt\}$
- $S^{(3)}(\omega) = \{\int_0^T d\omega_{t_1} d\omega_{t_2} d\omega_{t_3}, \int_0^T dt_1 d\omega_{t_2}, \int_0^T d\omega_{t_1} dt_2\}$
- ...

# Cubatures on Wiener Space

The smallest non-trivial cubature has degree 3 and consists of two paths:



We call it the “Rademacher cubature”.

# Cubatures on Wiener Space

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It is known that cubatures of arbitrarily high degree exist [Lyons and Victoir, 2004], but they are difficult to construct. Known cubatures:

- *Degree 3*: Rademacher cubature
- *Degree 5*: [Lyons and Victoir, 2004]
- *Degree 7*: [Hayakawa and Tanaka, 2022, Malyarenko and Nohrouzian, 2022]
- *Degree 9+*: [Koepernik et al., 2026]<sup>1</sup> (our work)

The error of the cubature approximation to  $\mathbb{E}[f(y_h)]$  is of order  $O(h^{(D+1)/2})$  for small  $h$  [Lyons and Victoir, 2004]  $\rightarrow$  subdivide  $[0, T]$  into  $N$  subintervals and apply the cubature on each of them  $\rightarrow$  total number of summands  $M = |\text{cubature}|^N$ , and

$$\text{global error} = O\left(N^{-(D-1)/2}\right) = O\left((\log M)^{-(D-1)/2}\right) \gg O\left(1/\sqrt{M}\right)$$

Monte-Carlo Error

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<sup>1</sup>Degree 9 technically constructed by [Gyurkó and Lyons, 2011] but not practically usable

# Our Contributions

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Our contributions are threefold:

- (i) We introduce a scalable algorithm that automatically constructs cubature formulae of arbitrary degree. We obtain cubatures with **degree up to 19** within hours on modest hardware.
- (ii) We prove that the error of the cubature approximation goes to zero *globally* on  $[0, T]$  as the degree  $D \rightarrow \infty$  (under some assumptions).
- (iii) We demonstrate empirically that the cubature approximation error with our cubatures (applied globally) is of order  $O(1/M)$ .

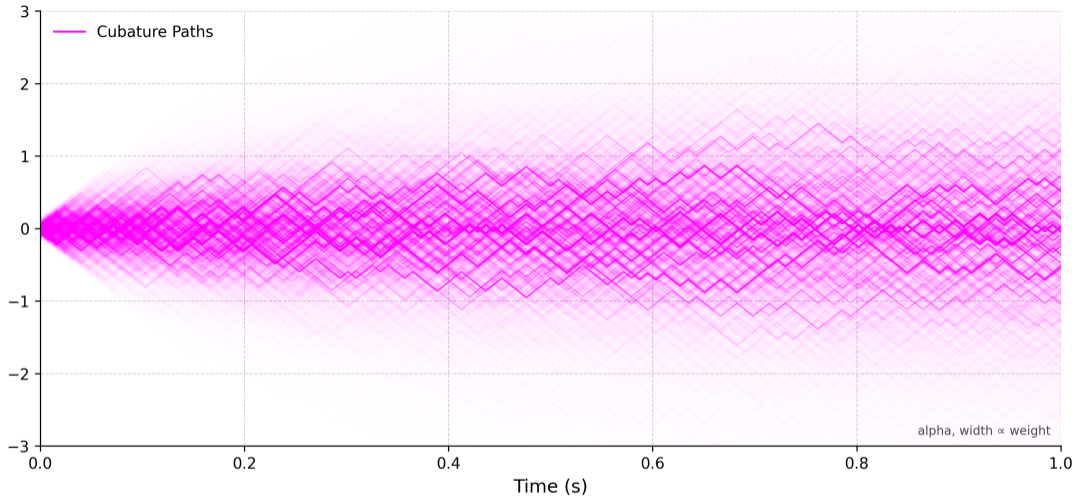


Figure: Degree 17 cubature, consisting of 3194 paths.

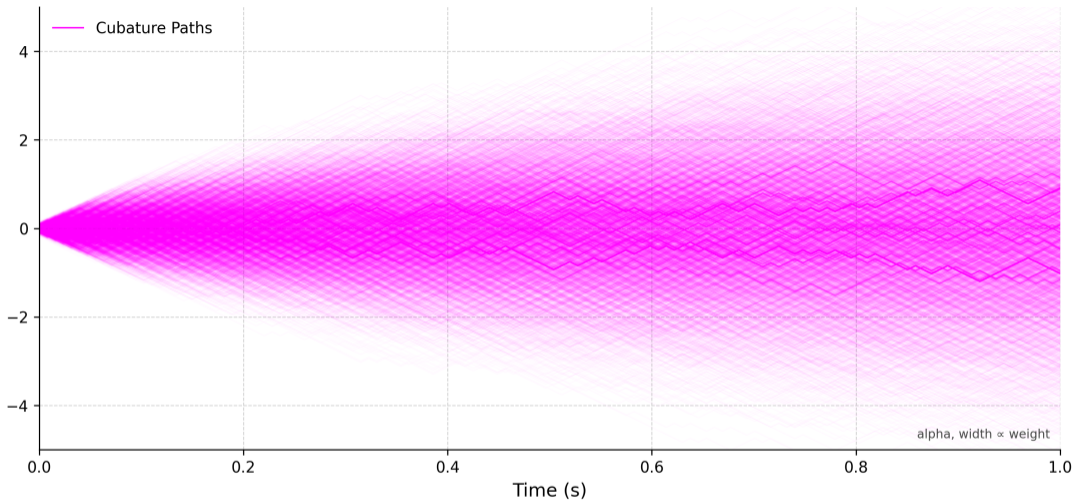


Figure: Degree 19 cubature, consisting of 8362 paths.

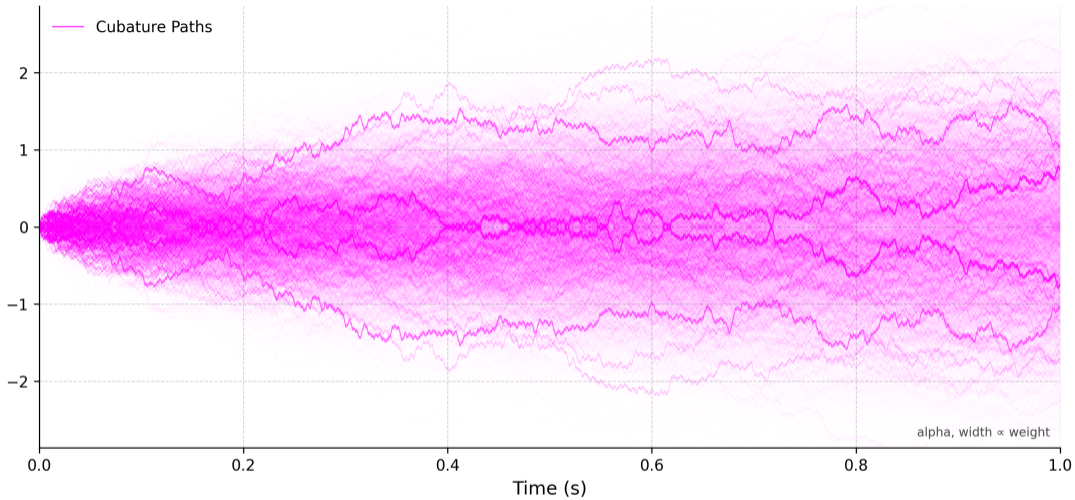
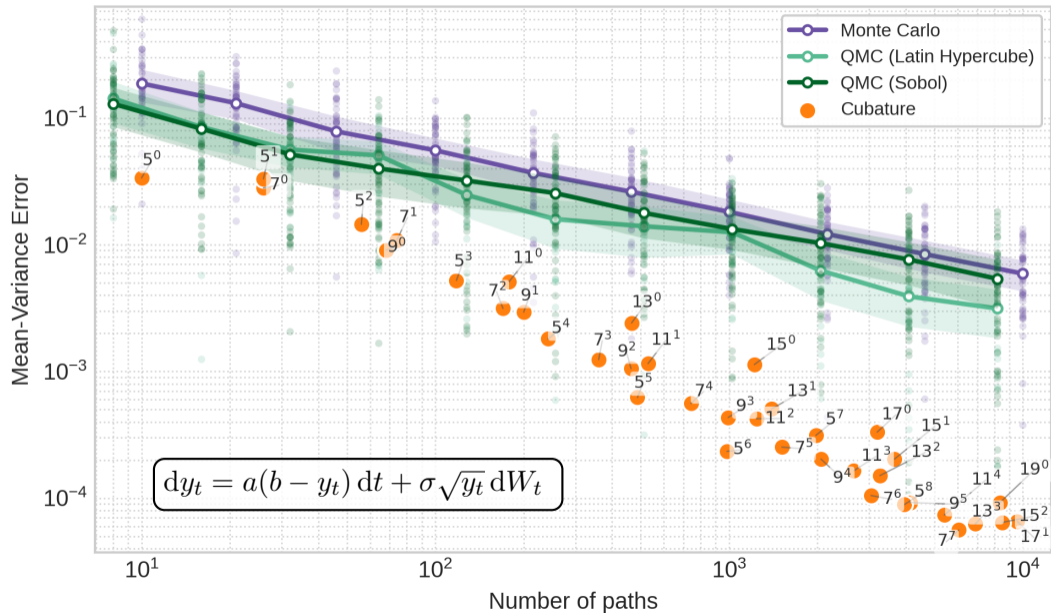
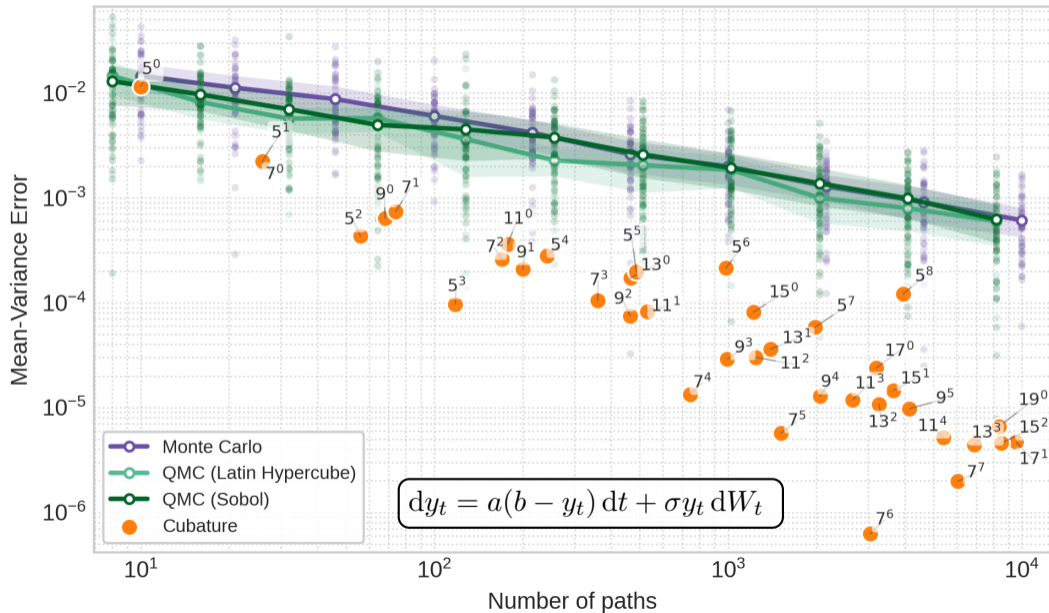


Figure: Degree 7 cubature (with “dyadic depth” 7).

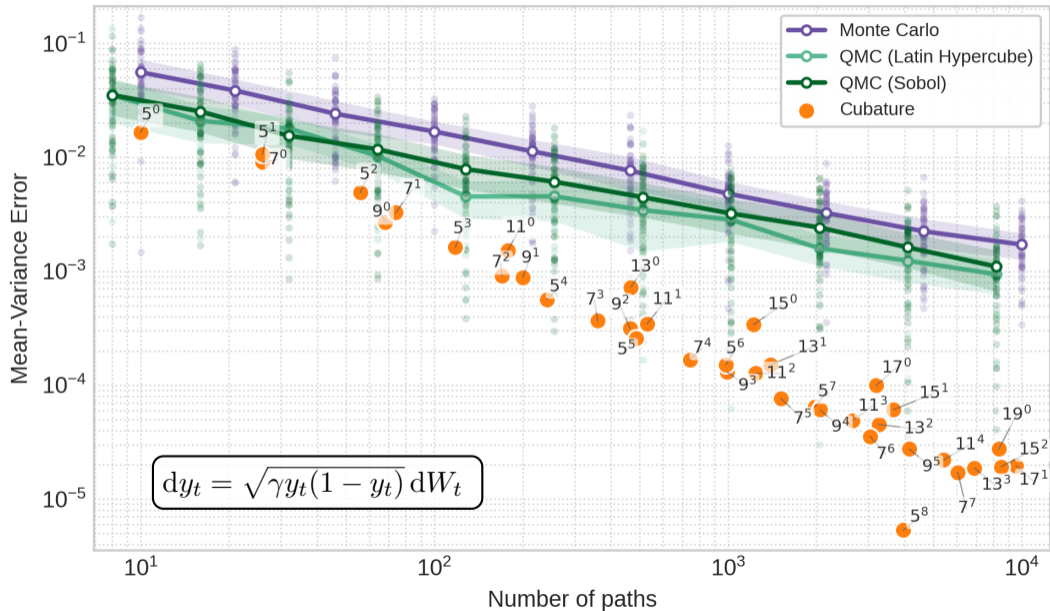
# Cox-Ingersoll-Ross (CIR) Model



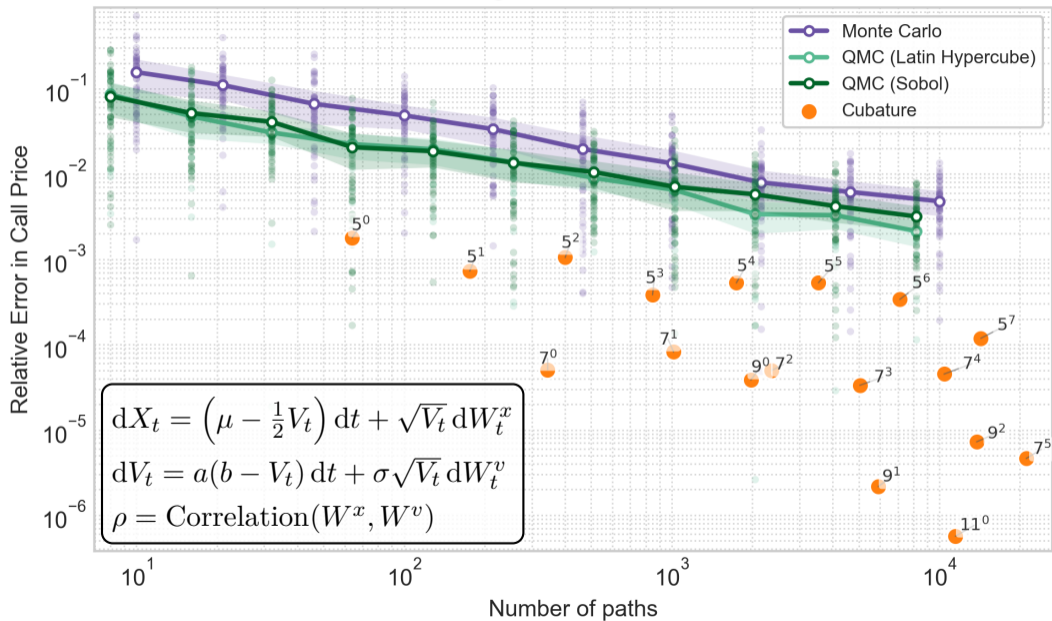
# Inhomogeneous Geometric Brownian Motion (IGBM)

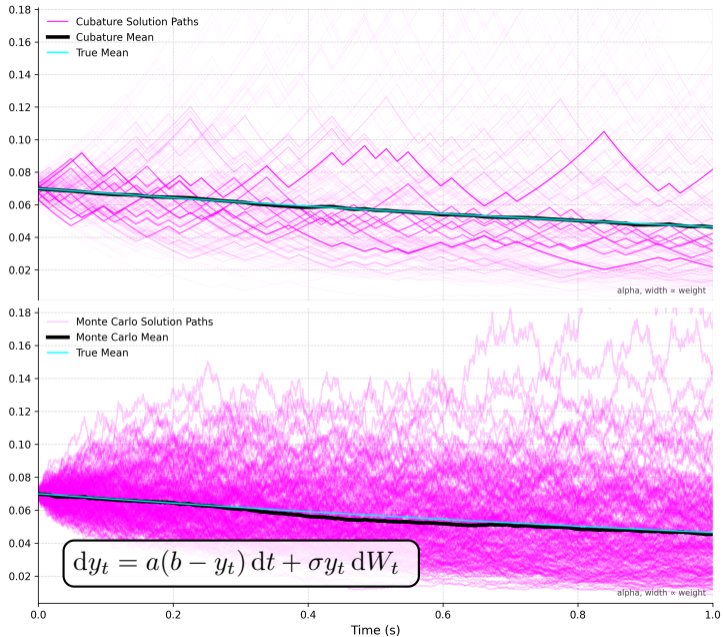


# Wright-Fisher Diffusion ( $t_1 = 1.0$ )



# Log-Heston Model





## Step 1: Construction of Approximate Cubature

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A discrete random walk on  $[0, T]$  with  $N$  steps of sizes  $\pm\sqrt{T/N}$  is an approximate cubature of size  $2^N$  with error  $O(1/N)$  (for fixed degree  $D$ ).

**Key observation:** The expected signature of this random walk does not change up to degree  $D$  if its  $N$  increments are only  $D$ -wise independent  $\rightarrow$  binary orthogonal arrays.

**Outcome:** Approximate cubature of size  $N^{\lfloor D/2 \rfloor}$  with error  $O(1/N)$ .

## Step 2: Recombination

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*Recombination* is an algorithm that takes a set of  $M$  points  $x_i \in \mathbb{R}^m$  with probability weights  $(\lambda_i)_{i=1}^M$ , and returns a subset  $(x_i)_{i \in I}$  of points with probability weights  $\bar{\lambda}_i$  of size  $|I| \leq m + 1$  such that

$$\sum_{i \in I} \bar{\lambda}_i x_i = \sum_{i=1}^M \lambda_i x_i.$$

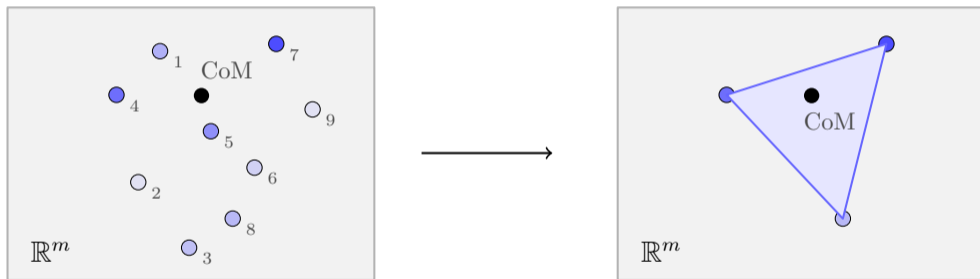
This algorithm is highly parallelisable and can be efficiently implemented on GPU. Thinking of the paths making up the approximate cubature as vectors of signature terms, we have

$$m(D) = |\{I \in \{0, 1\}^* : I \text{ has degree at most } D\}| \in [2^{(D+1)/2}, 2^{D+1}].$$

**Outcome:** Approximate cubature of size at most  $\underbrace{m(D) + 1}_{=: M(D)}$  with error  $O(1/N)$ .

## Step 2: Recombination

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## Step 3: Sharpening

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Finally, we attempt to modify the weights of the approximate cubature to make it an exact cubature. Abstractly, the task is “*given points  $x_i \in \mathbb{R}^m$  and a point  $p \in \mathbb{R}^m$ , does  $p$  lie in the convex hull of  $(x_i)_{i=1}^M$ , and if so what are the weights?*”. Let  $X \in \mathbb{R}^{m \times M}$  be the matrix with the  $x_i$  as columns.

- (i) **Solve  $\min_{\lambda} \|X\lambda - p\|$  without non-negativity restriction:** Every  $\lambda \in \mathbb{R}^M$  with  $\sum_{i=1}^M \lambda_i = 1$  that minimizes  $\|X\lambda - p\|$  can be written as  $\lambda_0 + Vz$  for some  $z \in \mathbb{R}^k$ , where  $k \in \mathbb{N}$ ,  $V \in \mathbb{R}^{M \times k}$ , and  $\lambda_0 \in \mathbb{R}^M$  can be computed efficiently using the SVD of  $X$ . If  $\min_{\lambda} \|X\lambda - p\| > 0$  then we failed, otherwise proceed to step (ii).
- (ii) **Find non-negative solution:** For example, find  $z^* \in \mathbb{R}^k$  with  $\lambda^* = \lambda_0 + Vz^* \geq 0$  by doing gradient descent on the “non-negativity loss”  $\ell(z) = \sum_{i=1}^M \max(0, -z_i)^2$ . We also implemented an alternating projections method – which was much faster.

**Outcome:** Exact cubature of size at most  $M(D)$ .

# Overview

Degree $D$	Minimum $N$ $N$	Step 1 Output Size		Steps 2 & 3 Output Size
		naive $2^N$	actual $N^{\lfloor D/2 \rfloor}$	$M(D)$
5	32	$4.3 \times 10^9$	$1.0 \times 10^3$	10
7	32	$4.3 \times 10^9$	$3.3 \times 10^4$	26
9	64	$1.8 \times 10^{19}$	$1.7 \times 10^7$	68
11	64	$1.8 \times 10^{19}$	$1.1 \times 10^9$	178
13	64	$1.8 \times 10^{19}$	$6.9 \times 10^{10}$	466
15	64	$1.8 \times 10^{19}$	$4.4 \times 10^{12}$	1220
17	64	$1.8 \times 10^{19}$	$2.8 \times 10^{14}$	3194
19	128?	$3.4 \times 10^{38}$	$9.2 \times 10^{18}$	8362

## Optional Step 4: Dyadic Construction

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Given a cubature  $(\lambda_i, \omega_i)_{i=1}^M$ ,

- (i) Set  $\tilde{\omega}_i(t) = \frac{1}{\sqrt{2}}\omega_i(2t)$  (rescales the paths to the interval  $[0, T/2]$ )
- (ii) Build the *product cubature*  $(\lambda_i \lambda_j, \tilde{\omega}_i \odot \tilde{\omega}_j)_{i,j=1}^M$ , where  $\odot$  denotes concatenation.  
→ It has size  $M^2$  and matches the expected Brownian signature on  $[0, T]$  as well as  $[0, T/2]$  and  $[T/2, T]$ .
- (iii) Apply recombination to the product cubature.

This can be repeated an arbitrary number of times. We call a cubature that matches the expected Brownian signature up to degree  $D$  on all dyadic intervals of lengths  $2^{-l}T$  for  $l \leq m$  a cubature with degree  $D$  and *dyadic depth*  $m$ .

# Theoretical Results

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## Definition

Let  $(\lambda_i^n, \omega_i^n)$  be a sequence of cubatures. Denote by  $y^n$  is the solution to an SDE driven by  $(\lambda_i^n, \omega_i^n)$ , and by  $y$  the solution to the same SDE driven by Brownian motion. Then we call the cubature(s) *asymptotically consistent* (for that SDE) if  $y^n \rightarrow y$  in distribution in the space of continuous paths.

We conjecture that this is true for the cubatures constructed by our algorithm under no assumptions other than that the degree  $D_n \rightarrow \infty$ .

# Theoretical Results

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## Theorem

*Suppose  $(\lambda_i^n, \omega_i^n)$  is a cubature constructed by our algorithm with degree  $D_n \rightarrow \infty$ , and dyadic depth  $m_n \in \mathbb{N}$  such that*

$$m_n \geq \left(\frac{3}{4} + \varepsilon\right) \log_2 N_n - C$$

*for some  $\varepsilon, C > 0$ . Then the cubature is asymptotically consistent.*

## Theorem

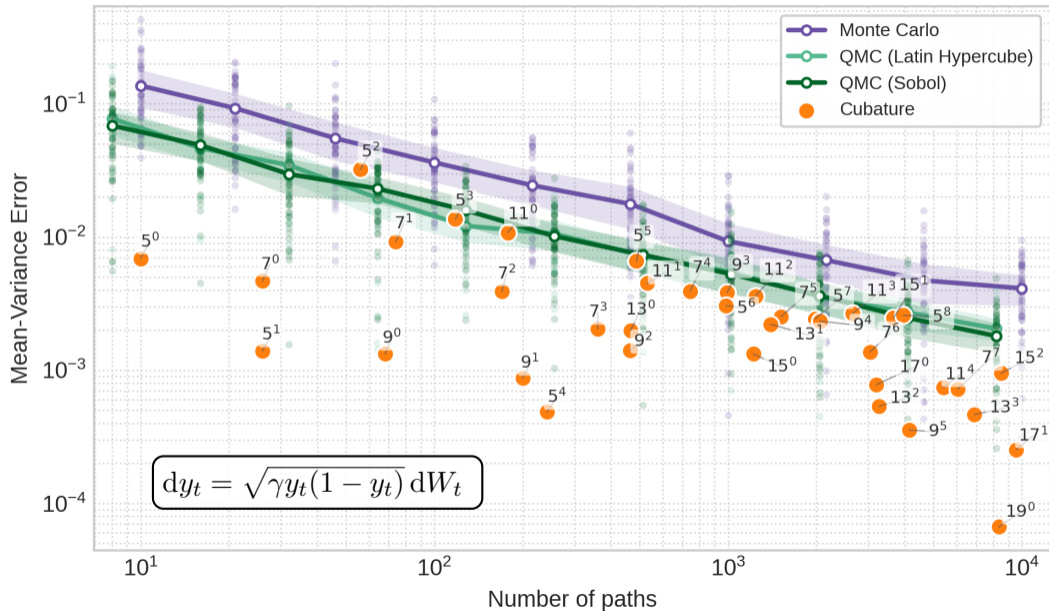
*Let  $(\lambda_i^n, \omega_i^n)$  be the approximate cubature resulting from step 1 of the algorithm, prior to recombination and sharpening, with  $N_n \rightarrow \infty$  and degree  $D_n \rightarrow \infty$ . Then it is asymptotically consistent.*

# Limitations

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- Cubature approximation can become inaccurate for non-smooth vector fields or very long time intervals.
- Algorithm technically works in any dimension but scales poorly, making it impractical for high-dimensional Brownian motion.  
Recombination scales  $\mathcal{O}\left(nm + \log_2\left(\frac{n}{2m}\right)m^3\right)$ , where  $n$  is the number of paths and  $m = m(D)$ .

# Wright-Fisher Diffusion ( $t_1 = 10.0$ )



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